



Radio Communications Research Group  
**UNIVERSITAT POLITÈCNICA DE CATALUNYA**

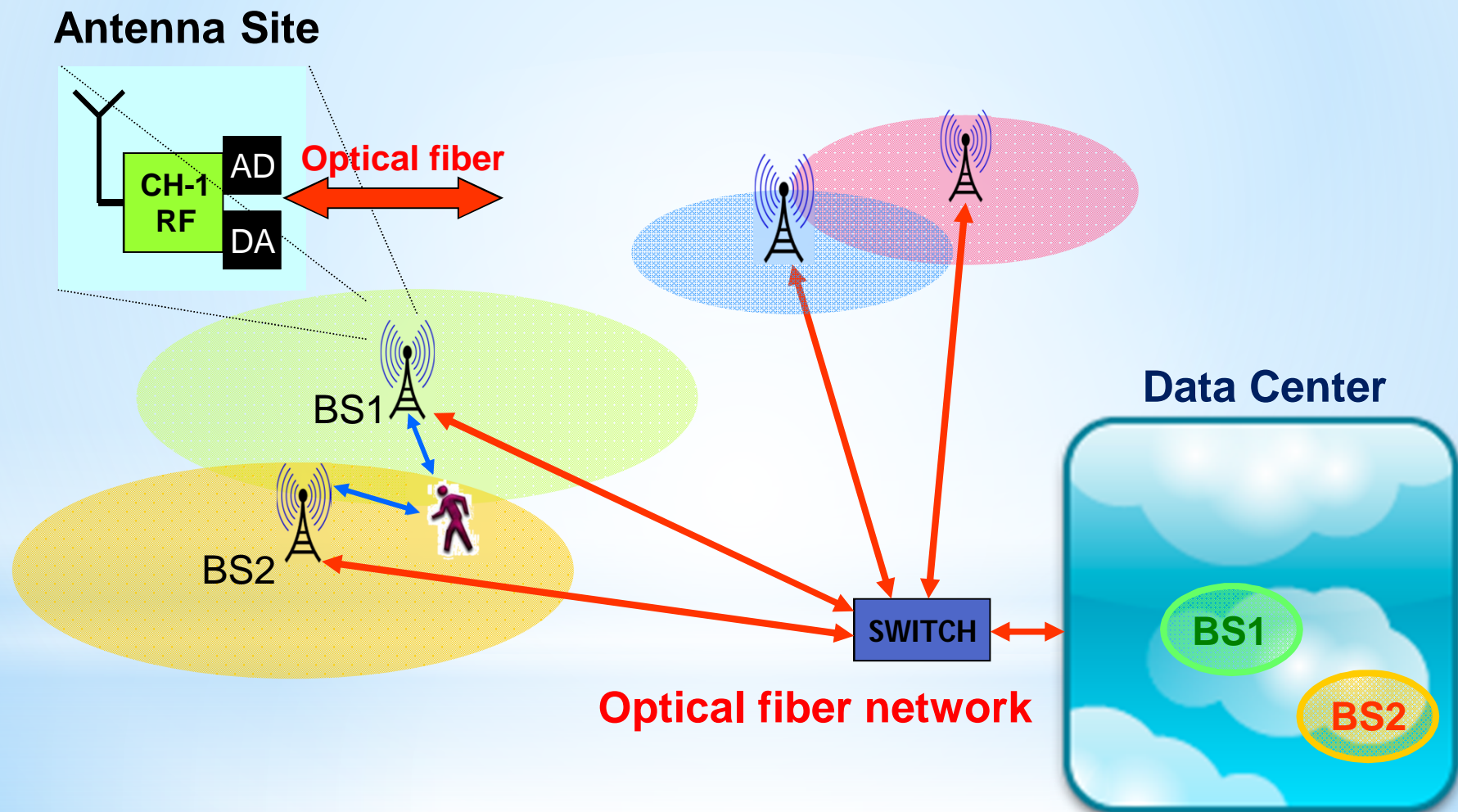
# **Evaluation of Computing Resource Management Methods for SDR Clouds**

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# SDR Cloud



# Real-Time Scheduling

- \* “In a real-time computing system *correctness* depends not only on the logical result of the computation but also on the time at which the results are produced”

J.A. Stankovic, A serious problem for next-generation systems.  
IEEE Computer, vol. 21, iss. 10, 1988.

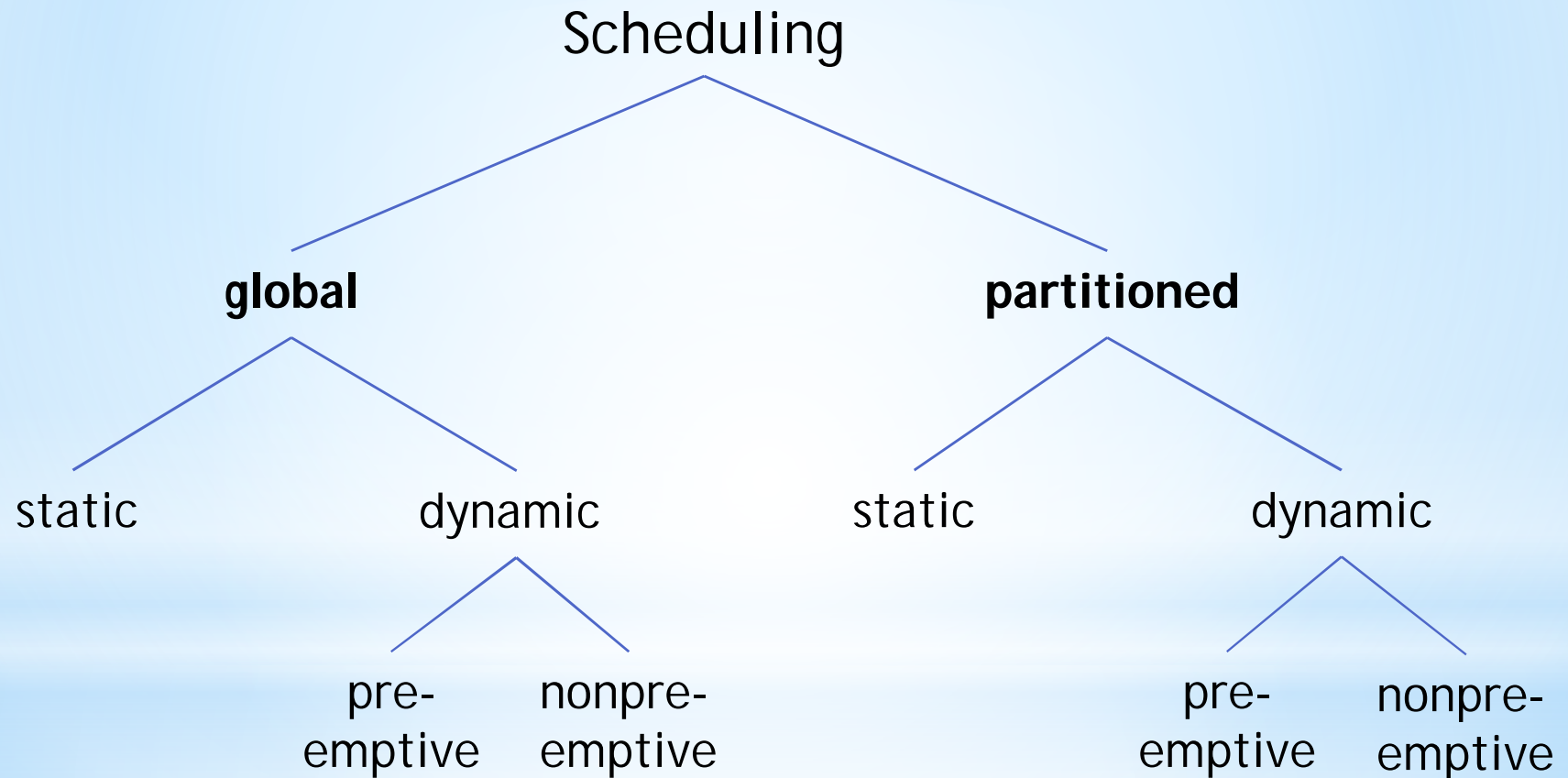
- \* SDR Cloud:
  - \* Allocate computing resources for real-time waveform execution (throughput and latency)
  - \* Runtime resource allocation (real-time)



# Problem Formulation

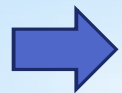
- \* Analyze different scheduling techniques for SDR clouds
  - \* in terms of complexity/resource overhead
  - \* *not* in terms of performance

# Scheduling Classification



# Global vs. Partitioned

Global	Partitioned
Global scheduling queue	Distributed schedulers
Implicit resource assignment and scheduling	Multiprocessor mapping, uniprocessor scheduling
Objective function can optimize for execution time (speedup)	Different objective for mapping and scheduling
Non-scalable	Scalable



Clustered or semi-partitioned scheduling

# Static vs. Dynamic

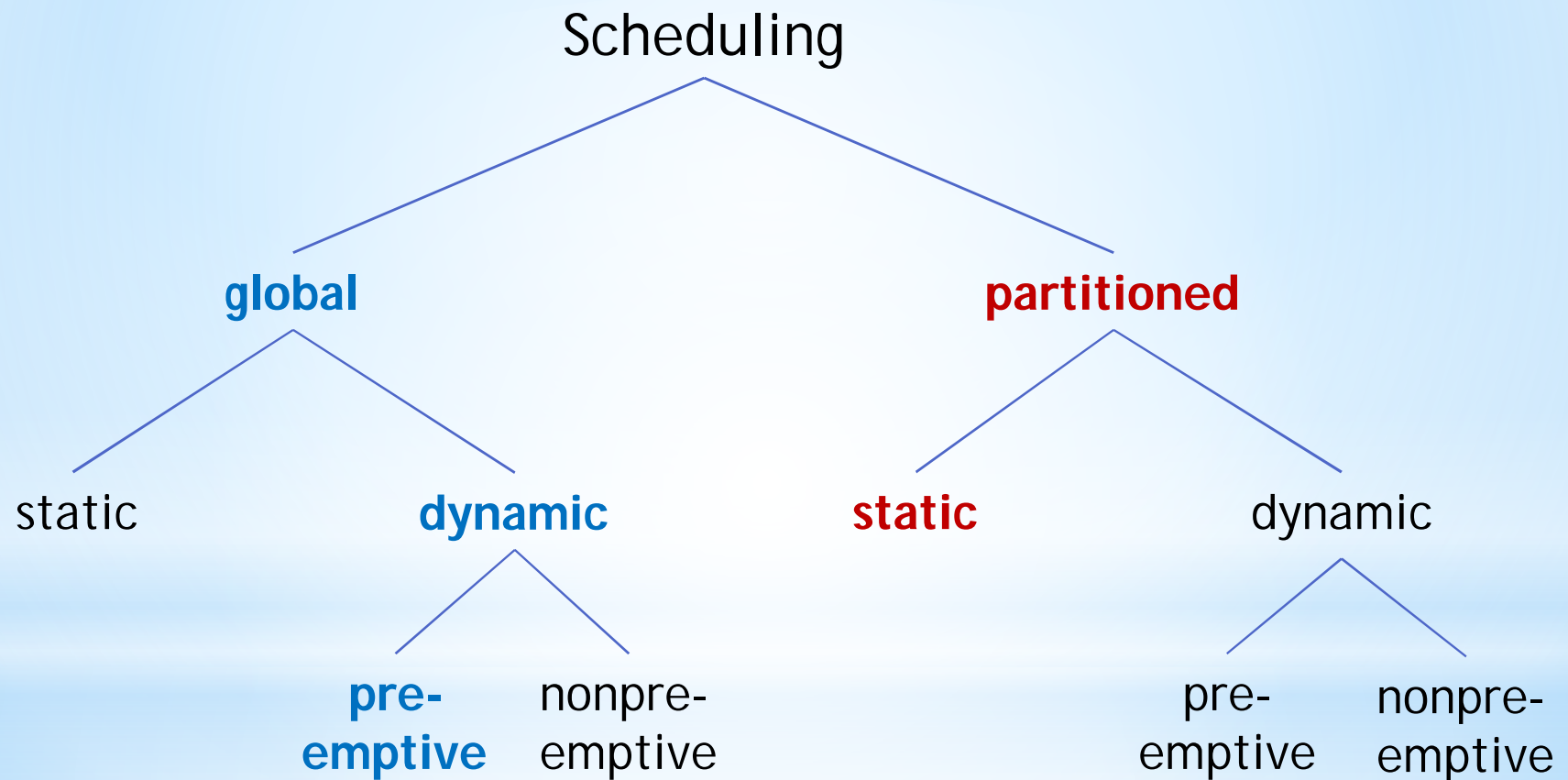
Static	Dynamic
Offline (compile time), before execution	Online (runtime), at execution
Deterministic performance	Nondeterministic performance
Avoid migrations → less overhead & fewer cache misses	Migrations → overhead & cache misses
Regular, periodic tasks with a priori information	Irregular, aperiodic tasks with unknown characteristics a priori
Runtime rescheduling costly	Easy to add new task at runtime

# Preemptive vs. Nonpreemptive

Preemptive	Nonpreemptive
Can stop an executing process and resume it later	Cannot stop a process in execution
Many context switches → overhead	Fewer context switches → less overhead
Dynamic task arrivals with different priorities	Periodic task arrivals or equal priorities
Simple schedulers	More complex schedulers



# Analysis



# Schedulers

## \* **Partitioned Static Scheduler (PSS)**

- \* Platform and waveform models (computing resources and requirements)
- \* Mapping of processing demands and data flows to processing and interprocessor bandwidth resources
- \* Local scheduling

## \* **Global Dynamic Preemptive Scheduler (GDPS)**

- \* Recalculates schedule of all active waveforms each time a user initiates or terminates a session

# Complexity Models

- \* Platform models
  - \* Cluster of  $n$  processing elements (part of SDR cloud data center)
- \* Waveform models
  - \* *Single-node waveform vs. processing chain of  $m$  tasks*
- \* Same scheduling complexity per waveform
- \* Different user arrival rates  $\lambda$
- \* Different processing loads  $\rho$  (# users  $u$ )

# Complexity Models

## \* Scheduling complexity per waveform

### \* Two models:

- *Single-node waveform:*  $t_1 = A \cdot n$
- *Distributed waveform processing:*  $t_2 = B \cdot m \cdot n^2$

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A, B: Scaling factors [s]

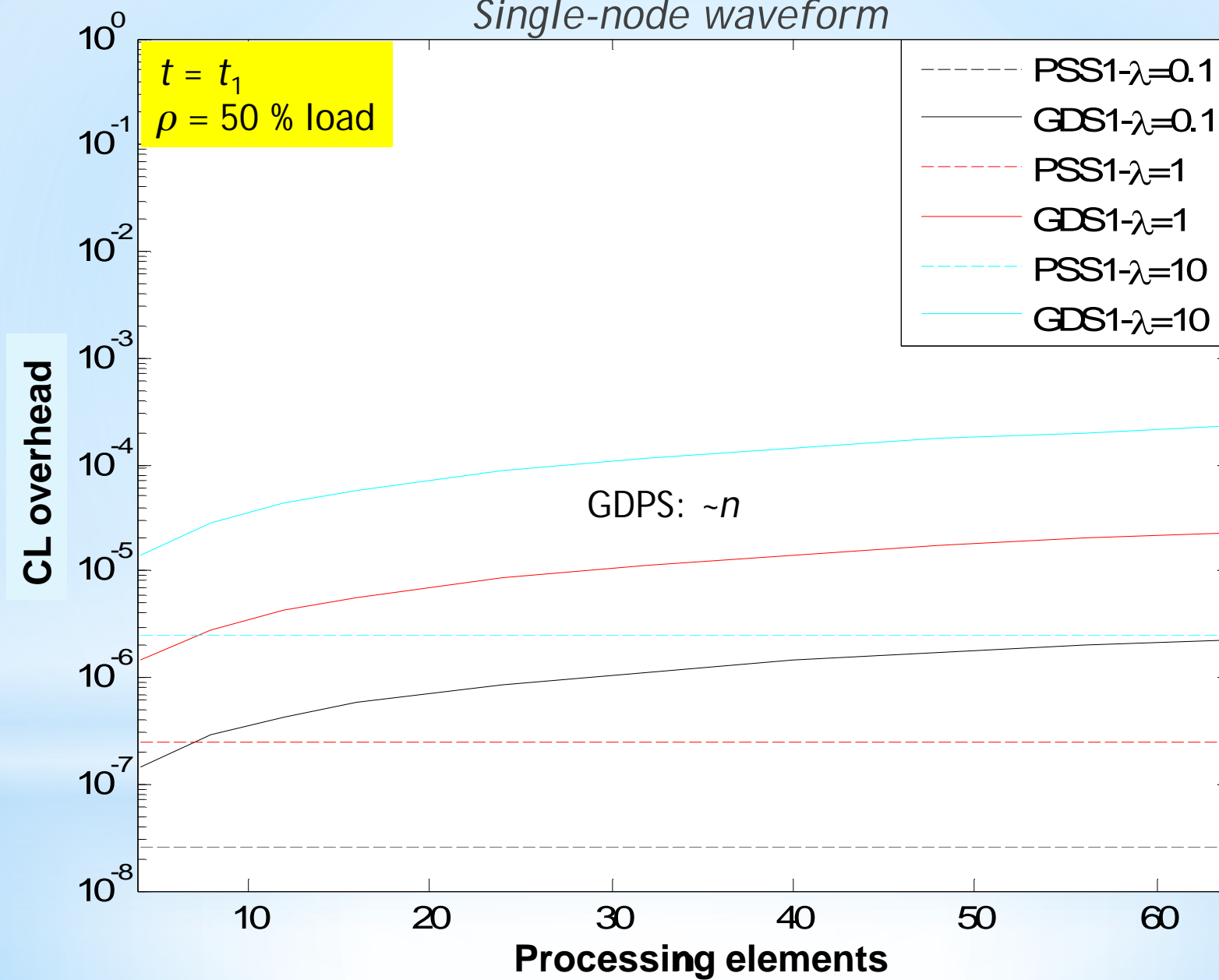
## \* Cluster processing overhead\*:

$$CL\text{-overhead}_{PSS} = \lambda \cdot t/n$$

$$CL\text{-overhead}_{GDPS} = 2\lambda \cdot u \cdot t/n \quad u \sim \rho \cdot n$$

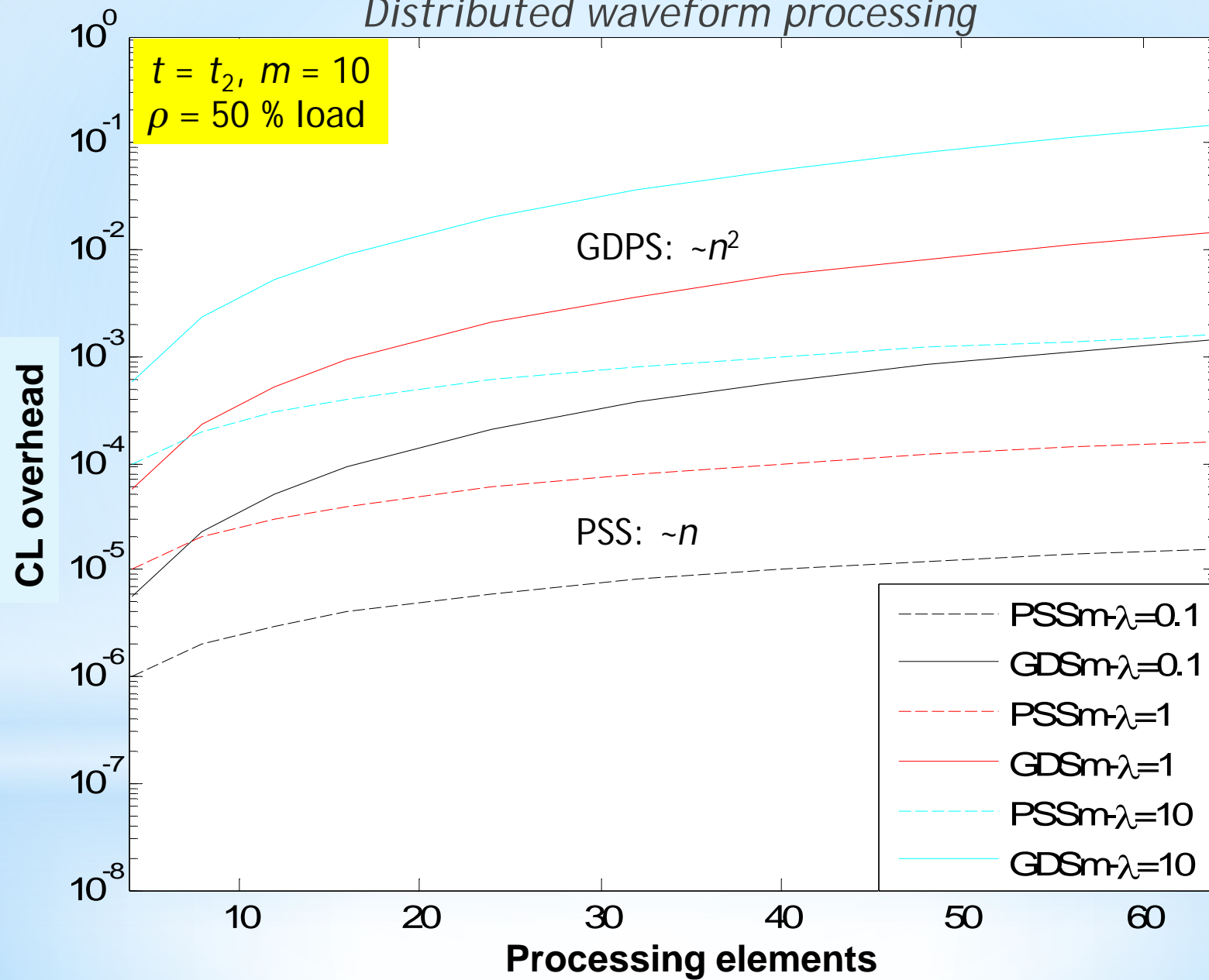
\*e.g. 0.1: 10 % cluster processing resources dedicated to scheduling

# Single-node waveform

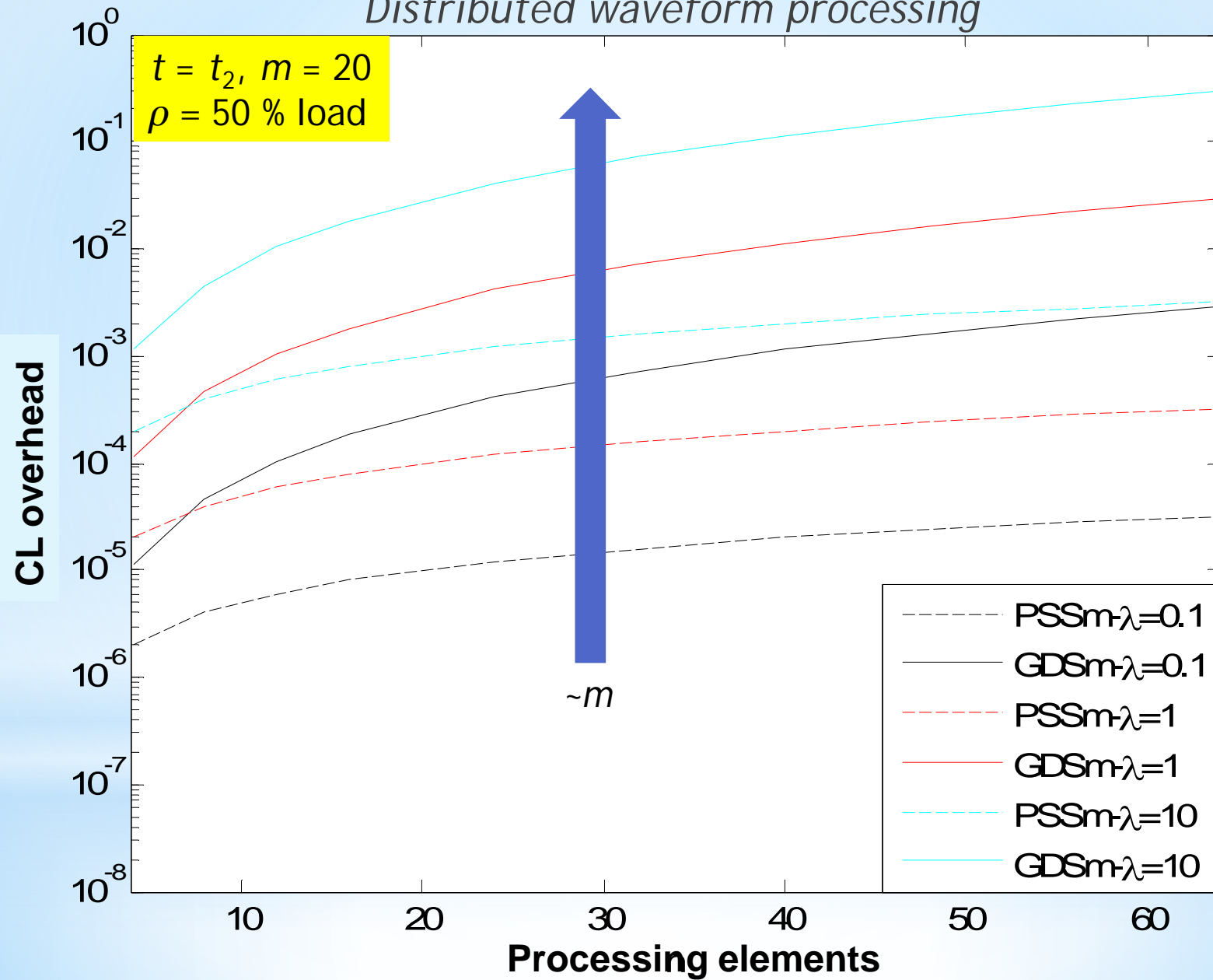




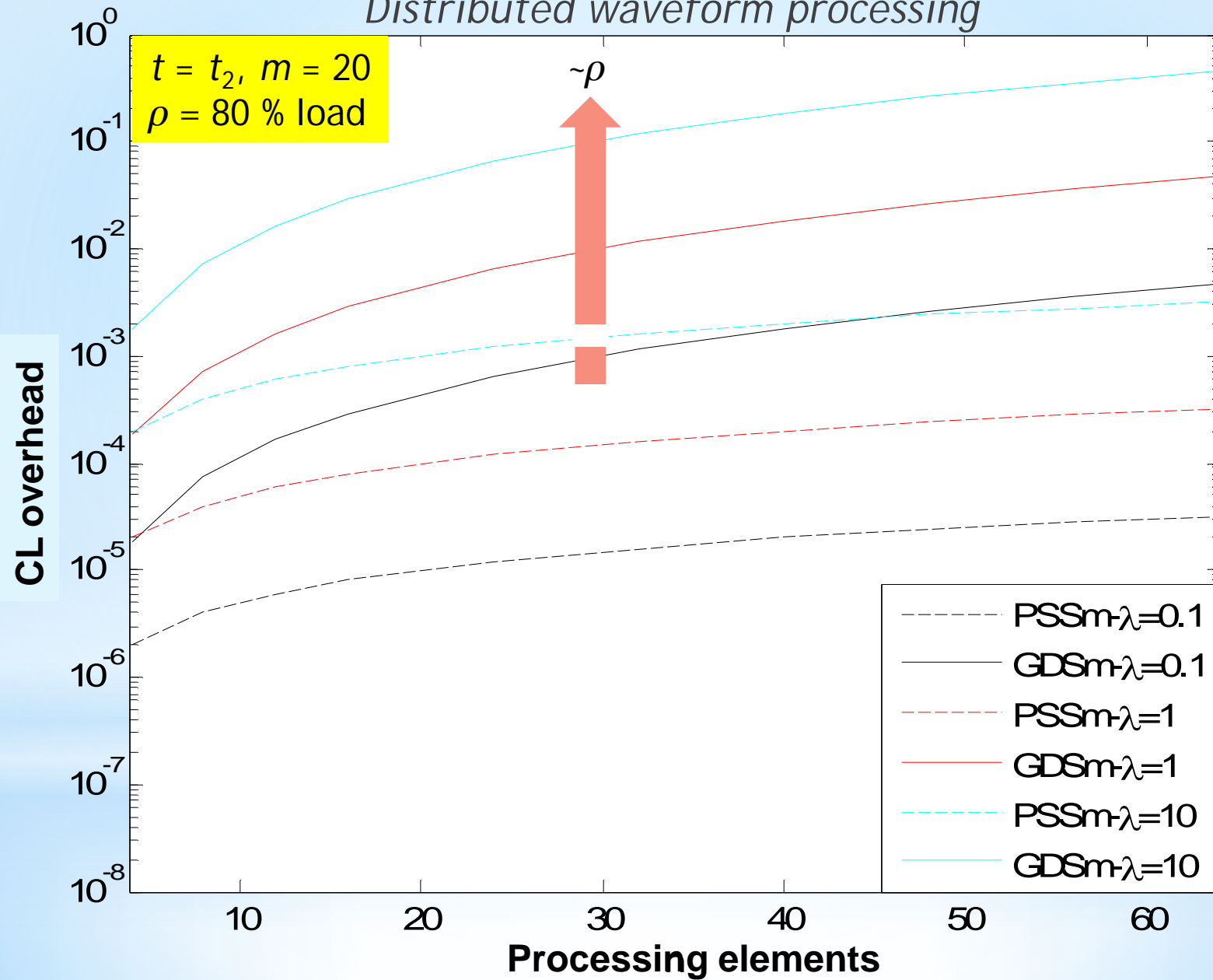
## Distributed waveform processing



## Distributed waveform processing



# Distributed waveform processing



# Conclusions

- \* Scheduling overhead increases
  - \* with the waveform granularity ( $m$ )
  - \* with the number of processing elements( $n$ )
  - \* inversely to the task execution time
- \* *Global-dynamic-preemptive* scheduling
  - + flexible
  - may incur significant resource overhead
- \* *Partitioned-static* scheduling scales better